## **Computational Challenges and Opportunities for Nuclear Astrophysics**



### **Bronson Messer**

Acting Group Leader Scientific Computing Group National Center for Computational Sciences

> Theoretical Astrophysics Group Oak Ridge National Laboratory

Department of Physics & Astronomy University of Tennessee





## Summary

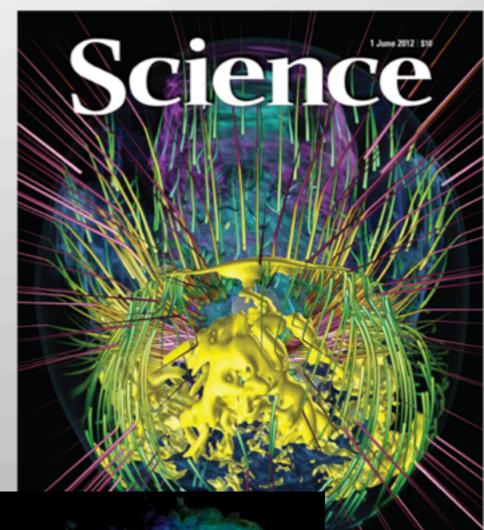
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- Stellar astrophysics is rife with unrealized parallelism, but architectural details and memory (i.e. cost, power) constraints will present considerable challenges. Additional support (for both "application scientists" and our CS/Math collaborators) will be required to surmount these challenges.
- Bulk-synchronous execution is a terrible way to try to exploit near-future architectures. A new programming model will require considerably more effort than a simple multi/many-core port.
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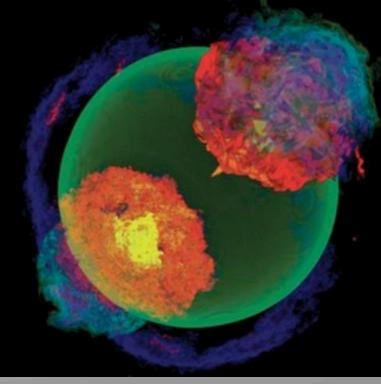


# Nuclear astrophysics INCITE allocations from 2010 - present

- Average number of cpu-hours/project -152 M
- In aggregate, just less than 10% of the total available each year from 2010 -2012
  - Allocations at NERSC are also aboveaverage in size
- This excellent record is due to
  - the formulation of large, important problems
  - demonstrated ability to efficiently exploit the largest computational platforms
- Will this trend continue to the exascale? Can we continue to solve big problems efficiently?



AAAS





Monday, July 23, 2012

#### The Effects of Moore's Law and Slacking <sup>1</sup> on Large Computations

Chris Gottbrath, Jeremy Bailin, Casey Meakin, Todd Thompson, J.J. Charfman

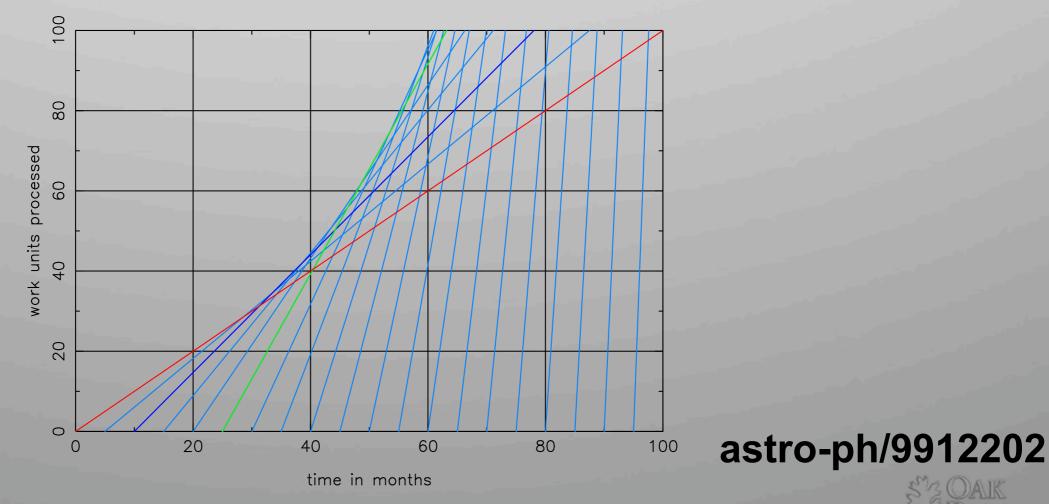
Steward Observatory, University of Arizona

<sup>1</sup>This paper took 2 days to write

#### Abstract

We show that, in the context of Moore's Law, overall productivity can be increased for large enough computations by 'slacking' or waiting for some period of time before purchasing a computer and beginning the calculation.

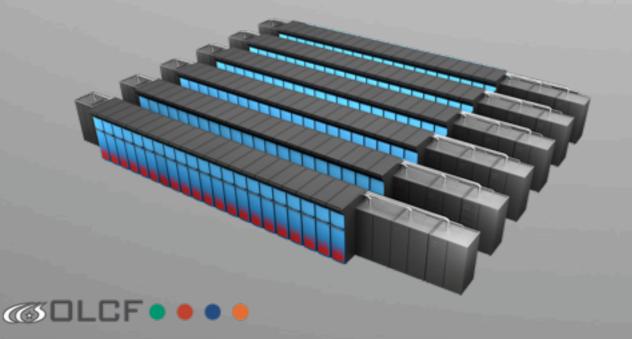
work and slack in the context of moores law





# **ORNL's "Titan" System**

- Upgrade of existing Jaguar Cray XT5
- Cray Linux Environment operating system
- Gemini interconnect
- 3-D Torus
- Globally addressable memory
- Advanced synchronization features
- AMD Opteron 6200 processor (Interlagos)
- New accelerated node design using NVIDIA multi-core accelerators
  - 2011: 960 NVIDIA M2090 "Fermi" GPUs ("titandev")
  - 2012: 20 PF NVIDIA "Kepler" GPUs
- 20 PFlops peak performance
  - Performance based on available funds
- 600 TB DDR3 memory (2x that of Jaguar)





Titan Specs	
Compute Nodes	18,688
Login & I/O Nodes	512
Memory per node	32 GB + 6 GB
NVIDIA "Fermi" (2011)	665 GFlops
# of Fermi chips	960
NVIDIA "Kepler" (2012)	>1 TFlops
Opteron	2.2 GHz
Opteron performance	141 GFlops
Total Opteron Flops	2.6 PFlops
Disk Bandwidth	~ 1 TB/s

### **Cray XK6 Compute Node ("titandev")**

XK6 Compute Node Characteristics AMD Opteron 6200 "Interlagos" **NVIDIA** 16 core processor @ 2.2GHz Cle Gen2 Tesla M2090 "Fermi" @ 665 GF with 6GB **NVIDIA** GDDR5 memory AMD Host Memory **T**3 32GB HT3 AMD 1600 MHz DDR3 AY Gemini High Speed Interconnect Upgradeable to NVIDIA's next generation "Kepler" processor in 2012 Four compute nodes per XK6 blade. 24 blades per rack **Multicore CPU** Many-Core GPU • 0 •



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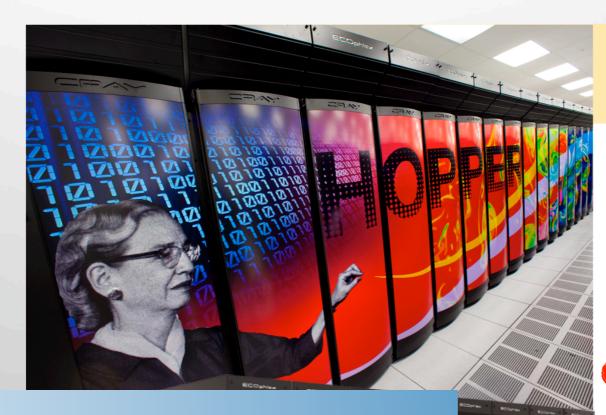
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NERSC

### This year



### NERSC-6 Grace "Hopper"

**Cray XE6** Performance 1.2 PF Peak 1.05 PF HPL (#5) Processor AMD MagnyCours 2.1 GHz 12-core 8.4.CFEOPs/core 24 cores/node 32-64 GB DDR3-1333 per node System Gemini Interconnect (3D torus) 6392 nodes 153,408 total cores **I/O** 2PB disk space 70GB/s peak I/O Bandwidth





### Franklin - Cray XT4

#### 38,288 compute cores

9,572 compute nodes

One quad-core AMD 2.3 GHz Opteron processors (Budapest) per node

4 processor cores per node

8 GB of memory per node

78 TB of aggregate memory

## 1.8 GB memory / core for applications

/scratch disk default quota of 750 GB



Light-weight Cray Linux operating system

No runtime dynamic, sharedobject libs

PGI, Cray, Pathscale, GNU compilers

### Last year

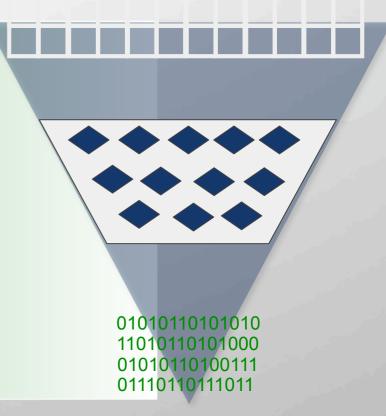


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## **Hierarchical Parallelism**

- MPI parallelism between nodes (or PGAS)
- On-node, SMP-like parallelism via threads (or subcommunicators, or...)
- Vector parallelism
  - SSE/AVX/etc on CPUs
  - GPU threaded parallelism



- Exposure of unrealized parallelism is essential to exploit <u>all</u> near-future architectures.
- Uncovering unrealized parallelism and improving data locality improves the performance of even CPU-only code.
- Experience with vanguard codes at OLCF suggests 1-2 personyears is required to "port" extant codes to GPU platforms.
  - Likely less if begun today, due to better tools/compilers

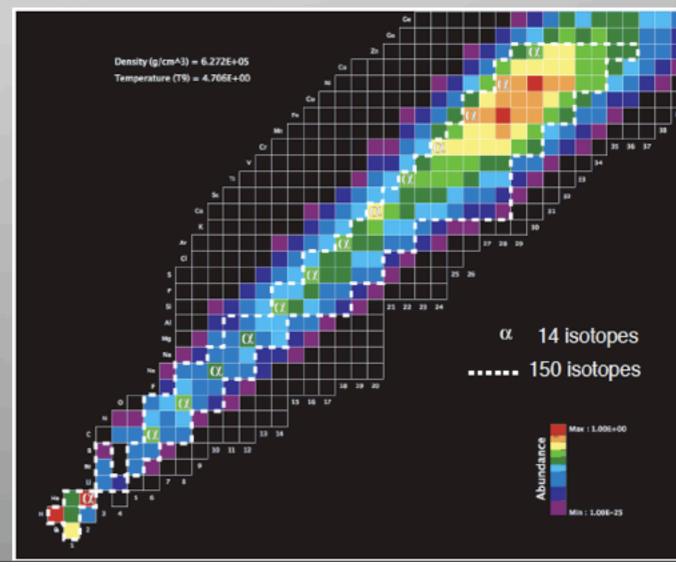


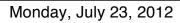
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((())) OLCF • • • • •

# Good news! Stellar astrophysics tends to have a lot of unrealized parallelism at present

- Simulation codes for stellar evolution and explosions
  - Exemplars of "multiphysics application codes"
  - Typically many degrees-of-freedom per spatial grid point
    - o radiation transport
    - o nuclear burning
  - Spatial domains typically parallelized via domain decomposition
- Related, computationallyintensive topics will, perhaps, have to work harder to identify additional parallelism outside of large stellar simulations, but plenty of opportunity exists.
  - high-density physics
  - nuclear structure and reactions





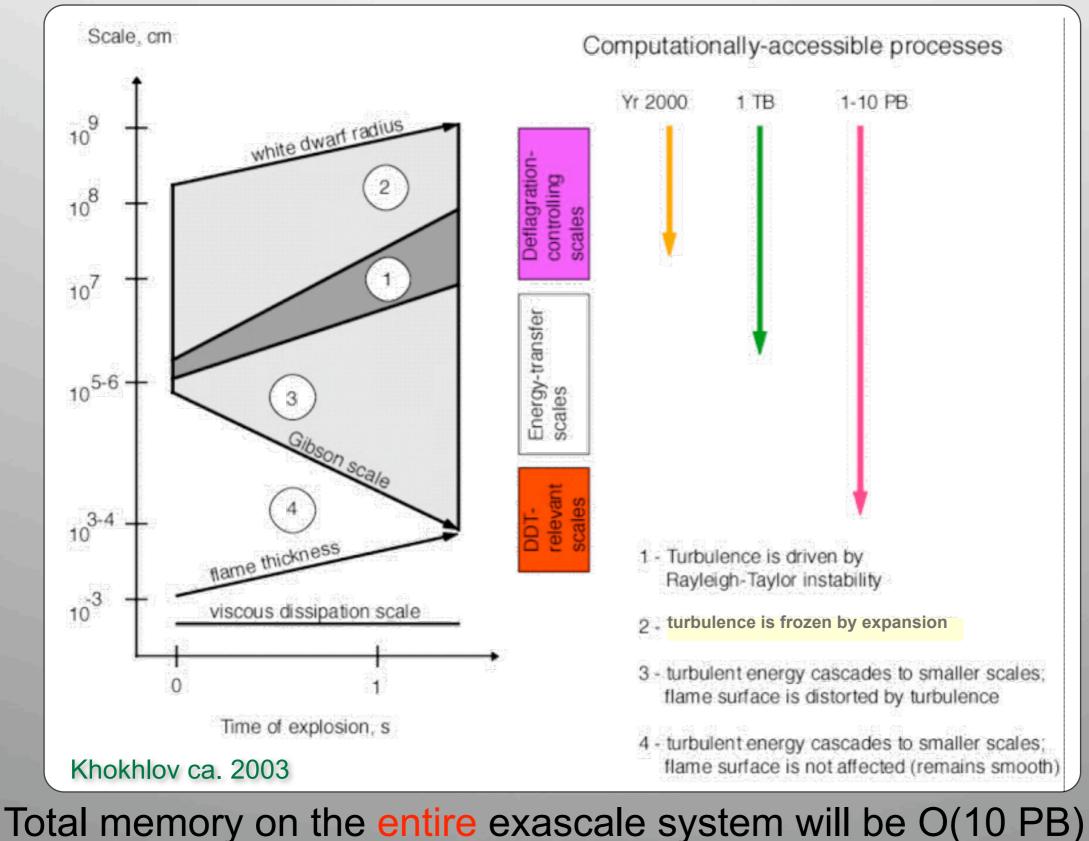
## **Posited Exascale Specs**

System attributes	2010	"2015"		"2018"	
System peak	2 PF	200 PF/s		1 Exaflop/s	
Power	6 MW	15 MW		20 MW	
System memory	0.3 PB	5 PB		32–64 PB	
Node performance	125 GF	0.5 TF	7 TF	1 TF	10 TF
Node memory BW	25 GB/s	0.1 TB/s	1 TB/s	0.4 TB/s	4 TB/s
Node concurrency	12	O(100)	O(1,000)	O(1,000)	O(10,000)
System size (nodes)	18,700	50,000	5,000	1,000,000	100,000
Total node interconnect BW	1.5 GB/s	150 GB/s	1 TB/s	250 GB/s	2 TB/s
MTTI	day	O(1 day)		O(1 day)	





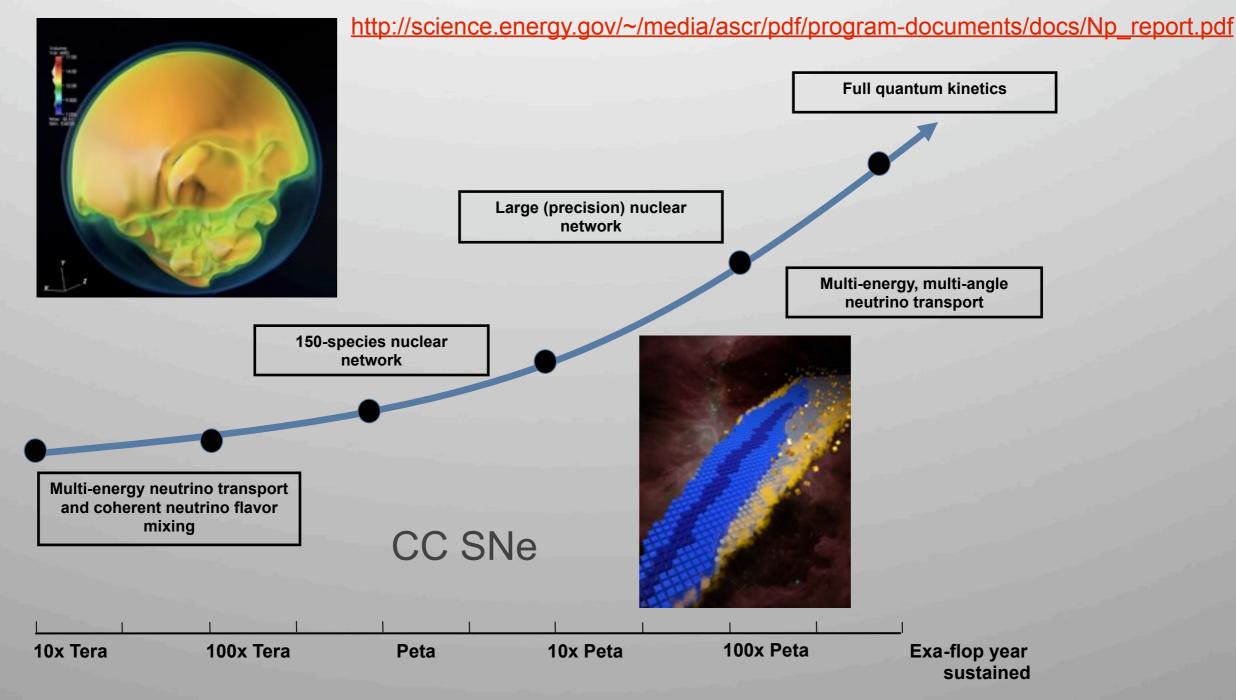
# Achieving high spatial (or phase-space, etc.) resolution will be very difficult.





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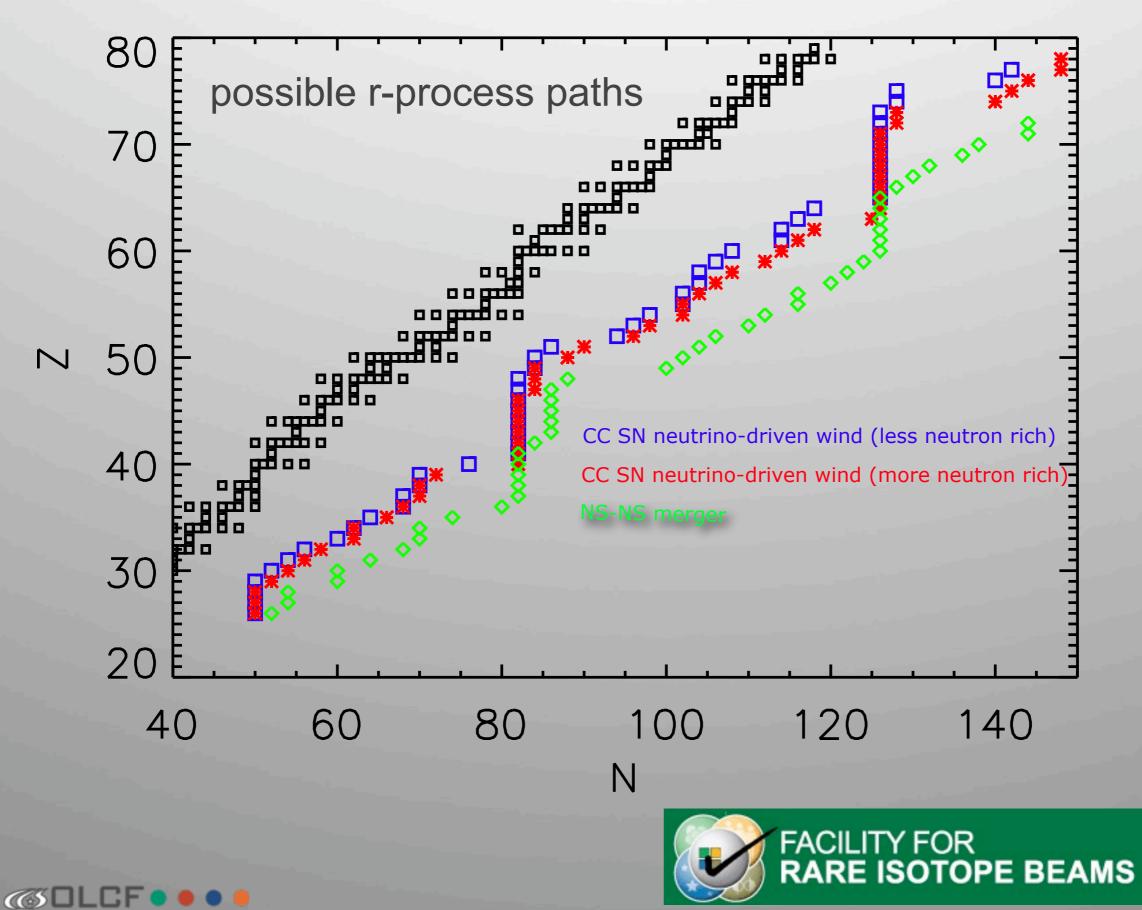
## **Example goals/highlights from NP Exascale** Workshop report (2009)



All of these goals are attainable, but will require new algorithms and implementations to bridge the gap to the posited architectures.



## Simulation is required to guide experiment





Locally bulk-synchronous programming model is not a viable path for maximum performance on these new platforms

- FLOP/s are cheap and moving data is expensive
- Even perfect knowledge of resource capabilities at every moment and perfect load balancers will not rescue billion-thread SPMD implementations of PDE simulations
- Cost of rebalancing frequently is too large, but the Amdahl penalty of failing to rebalance is fatal
- To take full advantage of asynchronous algorithms, we need to develop greater expressiveness in scientific programming
  - Create separate threads for logically separate tasks, whose priority is a function of algorithmic state, not unlike the way a time-sharing OS works
  - Join priority threads in a directed acyclic graph (DAG), a task graph showing the flow of input dependencies; fill idleness with noncritical work or steal work

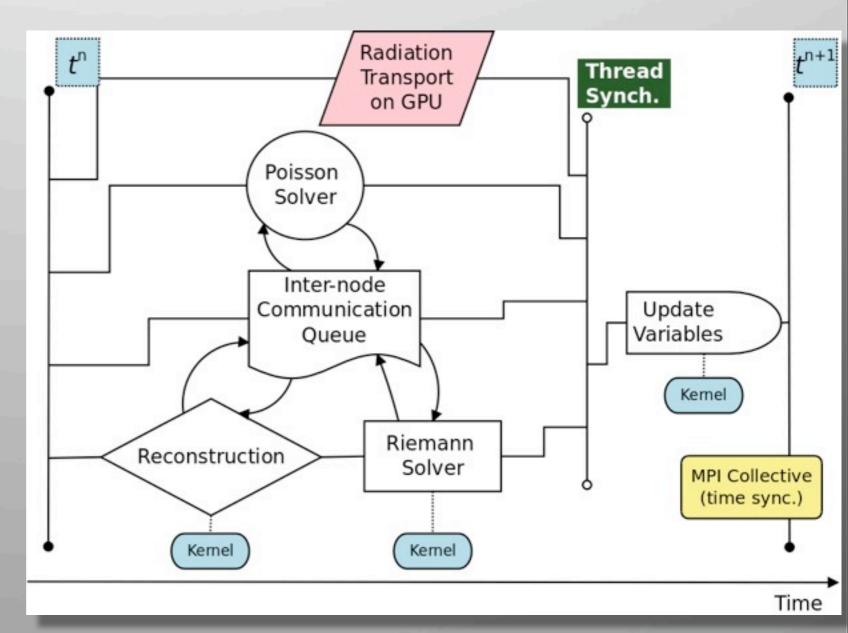
Comments taken directly from keynote address by David Keyes at EU-US HPC Summer School, June 2012





# Asynchronous execution models via task scheduling

- Examples exist already in other domains
  - MAGMA (linear algebra)
  - MADNESS (DFT)
  - Uintah (terrestrial combustion)
- Operator-split physics modules become "tasks" associated with execution threads







# Will the exascale (or before) machine be primarily a "strong-scaling" platform?

- Memory constraints provide a hard ceiling for spatial resolution and number of unknowns.
  - bytes/FLOP goes down by an order of magnitude
- Simulations will be certainly be larger, but likely not as large as one would expect if scaling with FLOPs is assumed.
  - no more than ~10x the number of MPI ranks?
  - this connotes no more than factors of ~2 in resolution in each dimension for 3D
  - OK: considerable understanding can be realized by fully exploring parameter space.
    - progenitor mass, rotation, metallicity
    - transport approximations, additional physics

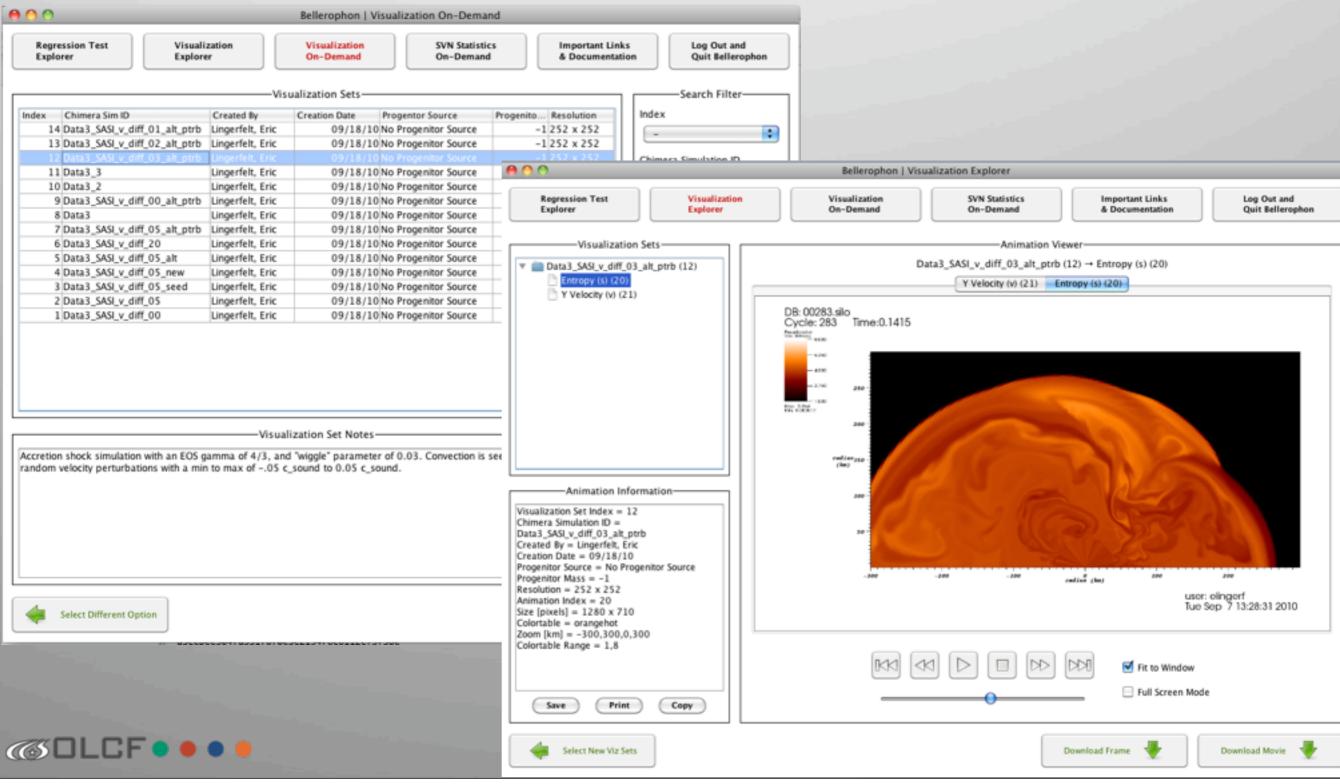




# Simulation, code, and data management become even harder



## Revision control, regression testing, viz, workflow...



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